

# Low complexity H.264 list decoder for enhanced quality real-time video over IP

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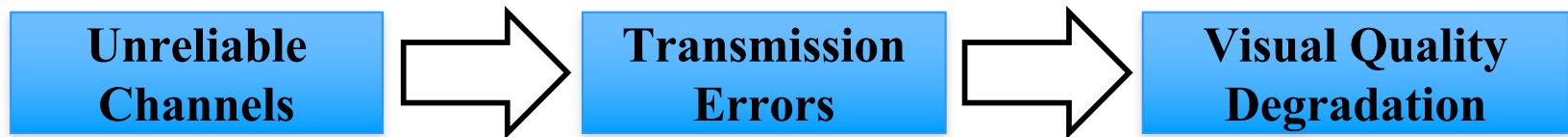
CCECE 2017

# Outline

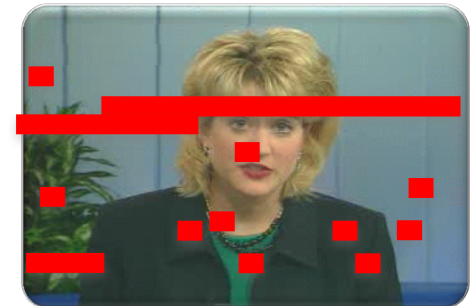
- ☐ Introduction
- ☐ Problem statement
- ☐ Literature review
- ☐ Proposed approach
- ☐ Experimental setup and results
- ☐ Conclusions



# Problem statement



- Corrupted Video Packets (bit errors)  
(e.g. attenuation, multi-path fading)
- Lost Video Packets  
(e.g. network congestion)



❖ **Compressed video streams are very vulnerable to errors**

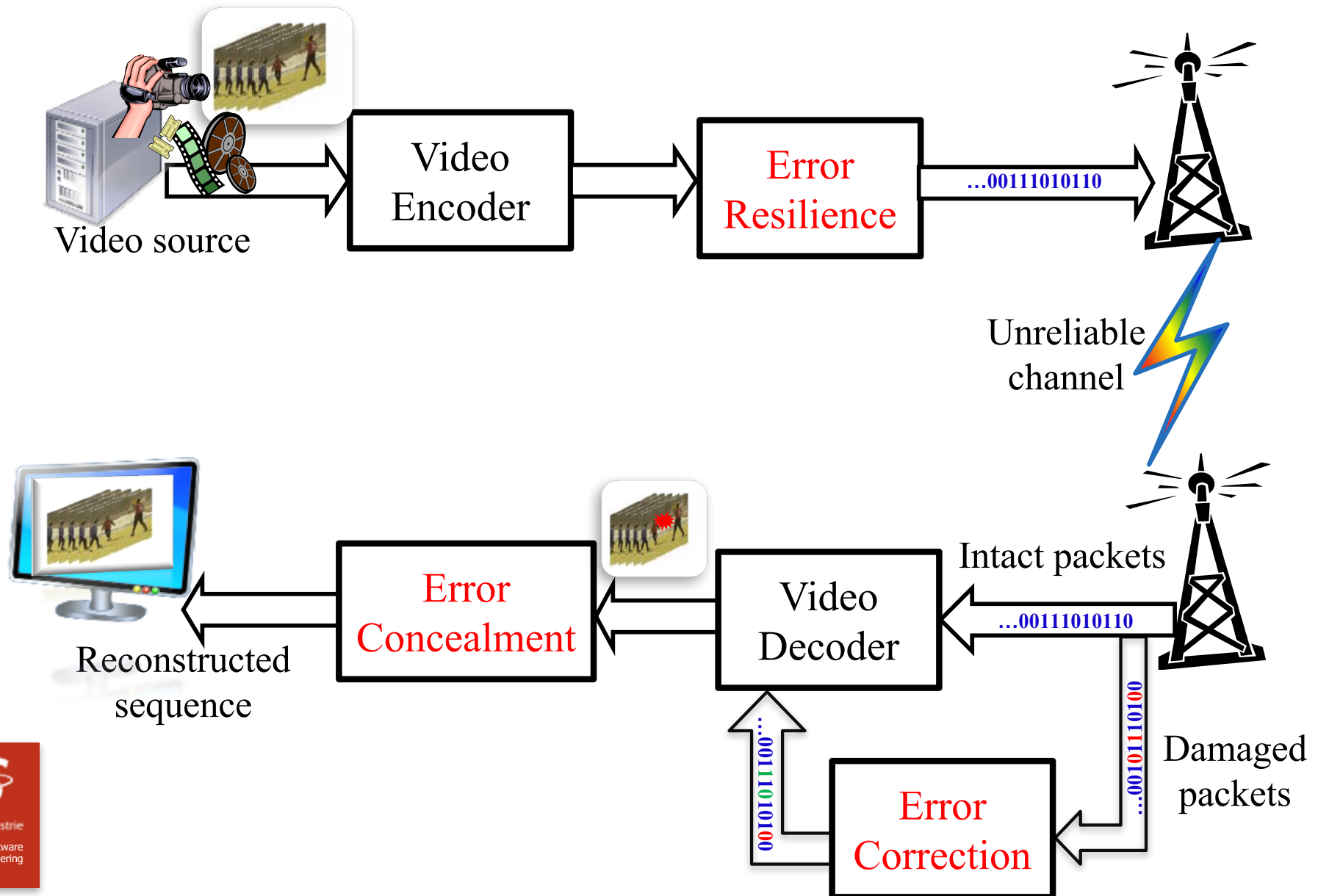


→  
Time

**Motion-compensated  
prediction process**

**Error propagation in compressed video!**

# Literature review



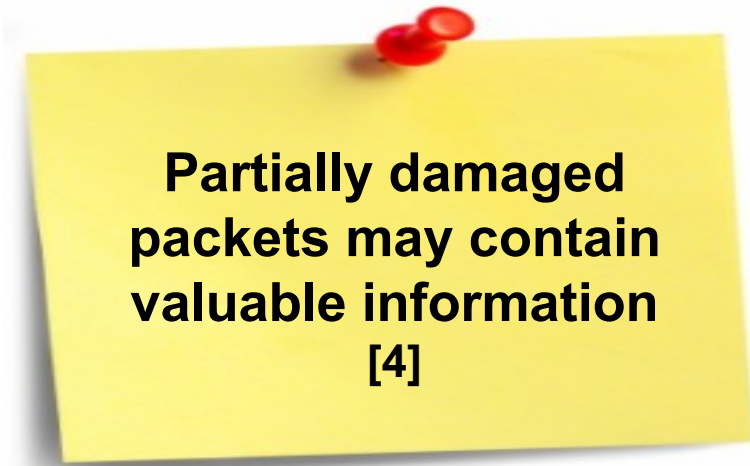
# Literature review



## ❑ Error Concealment:

Estimates missing pixels by using

- **Spatial correlation** [1]
- **Temporal correlation** [2]
- **Both spatial and temporal correlation**
  - - spatio-temporal boundary matching and partial differential equation (STBMA) [3]



**Partially damaged  
packets may contain  
valuable information  
[4]**

[1] J. Liu, et al., “Spatial error concealment with an adaptive linear predictor”, 2015.

[2] J. Zhou, et al., “Efficient motion vector interpolation for error concealment of H.264/AVC”, 2011.

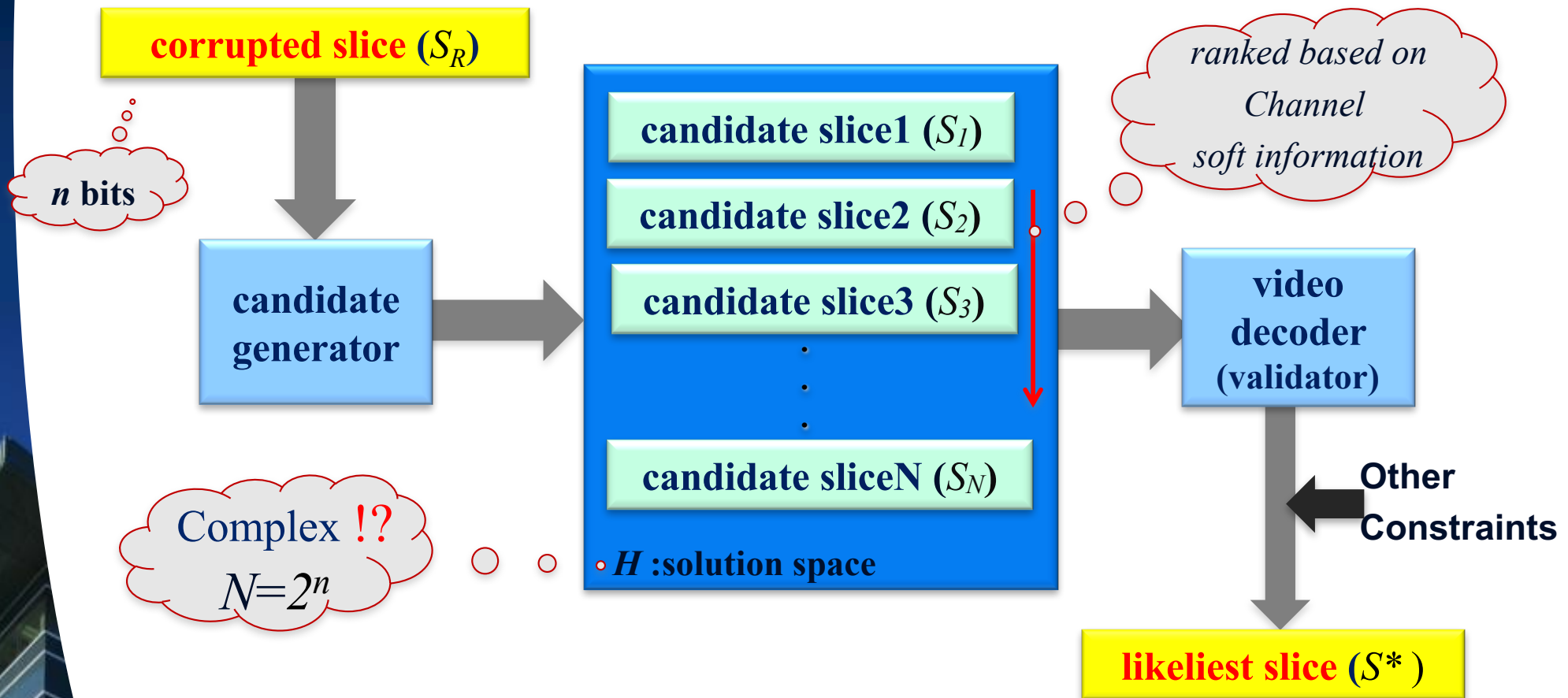
[3] Y. Chen, et al., “Video error concealment using spatio-temporal boundary matching and partial differential equation”, 2008.

[4] L. Superiori, et al., “Performance of a H.264/AVC error detection algorithm based on syntax analysis”, 2006.

# Literature review



## ❑ Error Correction: List decoding



- [5] D. Levine, et al., "Iterative joint source-channel Decoding of H.264 compressed video", 2007.
- [6] N. Nguyen, et al., "Iterative joint source- channel decoding for H.264 video transmission using virtual checking method at source decoder", 2010.
- [7] G. Sabeva, et al., "Robust decoding of H.264 encoded video transmitted over wireless channels", 2006.
- [8] C. Weidmann, et al., "Combined sequential decoding and error concealment of H.264 Video", 2004.
- [9] R. A. Farrugia , et al., "Robust decoder-based error control strategy for recovery of H.264/AVC video content", 2011.

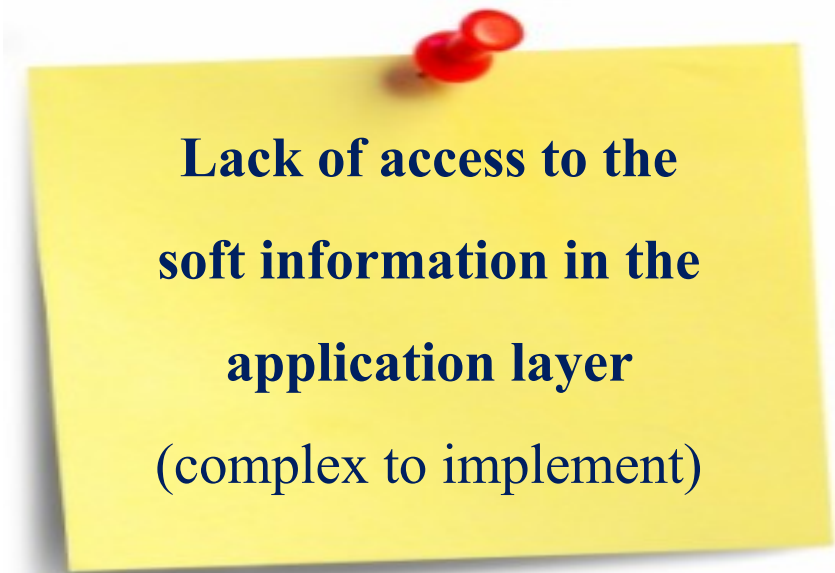
# Literature review



## ❑ Error Correction: Syntax Level

### Soft/Hard Output Maximum Likelihood Decoding (SO/HO-MLD) [10]

- Generate only a set of valid candidates at the correction process of each syntax element
- Any mistake in the decoding of a syntax element will propagate and reduce the performance



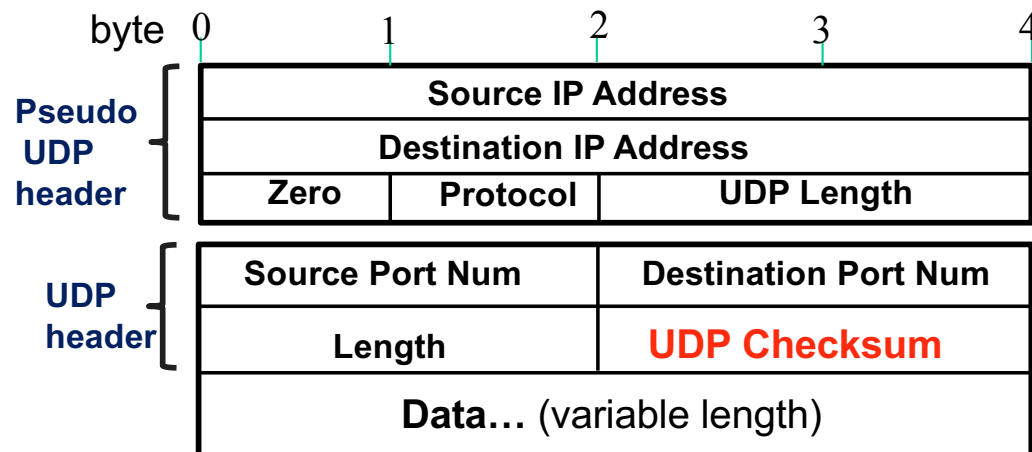
**Lack of access to the  
soft information in the  
application layer  
(complex to implement)**

# Proposed approach

- Reduction candidate list by analyzing the calculated UDP checksum value at the receiver side.

Checksum: **verify the integrity** of the delivered message

**UDP checksum:** one's complement of the one's complement sum of a *pseudo UDP header*, *UDP header* and application *data message*.





# Proposed approach

## TRANSMISSION

Bit Position	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Packet 1	0	0	1	0	0	1	1	0	1	1	0	1	0	0	0	1	1	1	0	0	1	1	0	0	1	0	0	1	1	0	0	0

Bit Position	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Second 16-bit word	0	0	1	0	0	1	1	0	1	1	0	1	0	0	0	1
First 16-bit word	1	1	0	0	1	1	0	0	1	0	0	1	1	0	0	0
Ones's compl. Sum	1	1	1	1	0	0	1	1	0	1	1	0	1	0	0	1
Checksum ( $C_T$ )	0	0	0	0	1	1	0	0	1	0	0	1	0	1	1	0

## UDP checksum calculation:

- Divide the packet into chunks of 16-bit words.
- Compute one's complement sum over all the words.
- Flip all the bits of the final sum (one's complement).

$$C_T = \sum_{i=0}^{m-1} \overline{W_i}$$

16-bit word

# Proposed approach

## RECEPTION

Bit Position	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Packet 1	0	0	1	0	0	1	1	0	1	1	0	1	0	0	0	1	1	1	0	0	1	1	0	0	1	0	0	1	1	0	0	0

Bit Position	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Second 16-bit word	0	0	1	0	0	1	1	0	1	1	0	1	0	0	0	1
First 16-bit word	1	1	0	0	1	1	0	0	1	0	0	1	1	0	0	0
Ones's compl. Sum	1	1	1	1	0	0	1	1	0	1	1	0	1	0	0	1
Checksum (C <sub>T</sub> )	0	0	0	0	1	1	0	0	1	0	0	1	0	1	1	0
	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
Checksum(C <sub>R</sub> )	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

The **zero** value of C<sub>R</sub> validates the received packet.

$$C_R = \sum_{i=0}^{m-1} \overline{W_i} + \overline{C_T}$$

If no error:

$$= \sum_{i=0}^{m-1} \overline{W_i} + \sum_{i=0}^{m-1} \overline{W_i} = \sum_{i=0}^{m-1} (W_i + \overline{W_i})$$

$$= \overline{FFFF} = 0000$$

the received versions of W<sub>i</sub> and C<sub>T</sub> are denoted as  $\overline{W_i}$  and  $\overline{C_T}$

# Proposed approach

## TRANSMISSION

Bit Position	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Packet 1	0	0	1	0	0	1	1	0	1	1	0	1	0	0	0	1	1	1	0	0	1	1	0	0	1	0	0	1	1	0	0	0

## RECEPTION

Corrupted

Bit Position	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Packet 1	0	0	1	0	0	0	1	0	1	1	0	1	0	0	0	1	1	1	0	0	1	1	0	0	1	0	0	1	1	0	0	0

Bit Position	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Second 16-bit word	0	0	1	0	0	0	1	0	1	1	0	1	0	0	0	1
First 16-bit word	1	1	0	0	1	1	0	0	1	0	0	1	1	0	0	0
Ones's compl. sum	1	1	1	1	0	1	1	1	0	1	1	0	1	0	0	1
$C_T$	0	0	0	0	1	1	0	0	1	0	0	1	0	1	1	0
	1	1	1	1	1	0	1	1	1	1	1	1	1	1	1	1
Checksum( $C_R$ )	0	0	0	0	0	1	0	0	0	0	0	0	0	0	0	0

Non-zero value of  $C_R$  : received packet is corrupted.

The checksum is indication the position of the error in modulo 16.



# Proposed approach

Error Type	$C_R$	$C_R$ Patterns
$1_j \longrightarrow 0_j$	<div>000 ... 0<b>1</b>0 ... 000</div> <div>Col. 15 j 0</div>	1000 0000 0000 0000 0100 0000 0000 0000 . 0000 0000 0000 0001
$0_j \longrightarrow 1_j$	<div>111 ... 1<b>0</b>1 ... 111</div>	0111 1111 1111 1111 1011 1111 1111 1111 : 1111 1111 1111 1110

The  $C_R$  pattern :

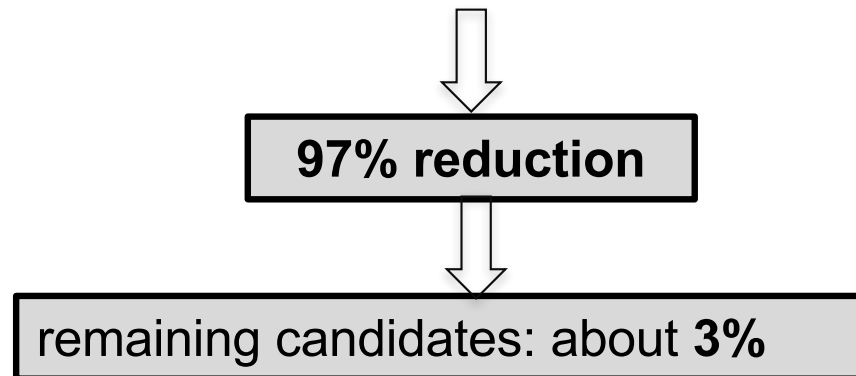
- potential error columns in the words
- type of the flipped bits



# Candidate Reduction

One bit in error in a packet of  $N$  bits:

- Traditional list decoding approaches:  $N$  candidates
- proposed approaches:  $N/32$  candidates.



Packet length (bits)	Average number of candidates		Eliminated candidates
	List Decoding	Proposed	
272	272	9	97%
768	768	24	97%
1120	1120	35	97%
5280	5280	165	97%

# Simulation setup

## ❖ Coding Parameters:

- **Baseline** profile of H.264 (simulator: **JM** software v.18.5)
- First 60 Frames of sequence: *Foreman* (352x288), *Crew* (704x576), *Opening Ceremony*, *Walk* (720x576).
- Coded in **IPPP...** format, Intra refresh rate of **30** frames, quantization parameters (QPs): 22, 27, 32, and 37
- Each slice (= a single row of MBs), encapsulated into RTP packets.
- UDP checksum calculated on RTP packet

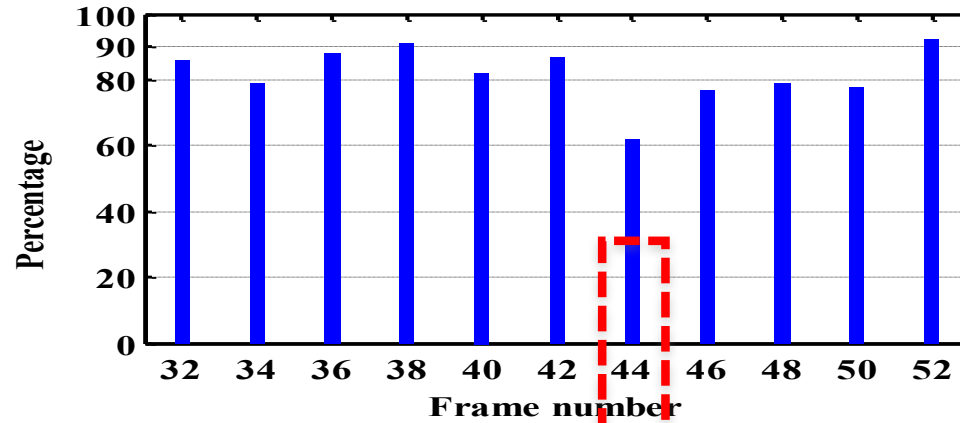


- ❖ Corrupted frame: **31<sup>th</sup>-59<sup>th</sup>** randomly
- ❖ Channel bit error rate (BER): [**10<sup>-7</sup>**, **10<sup>-6</sup>**]: one bit error

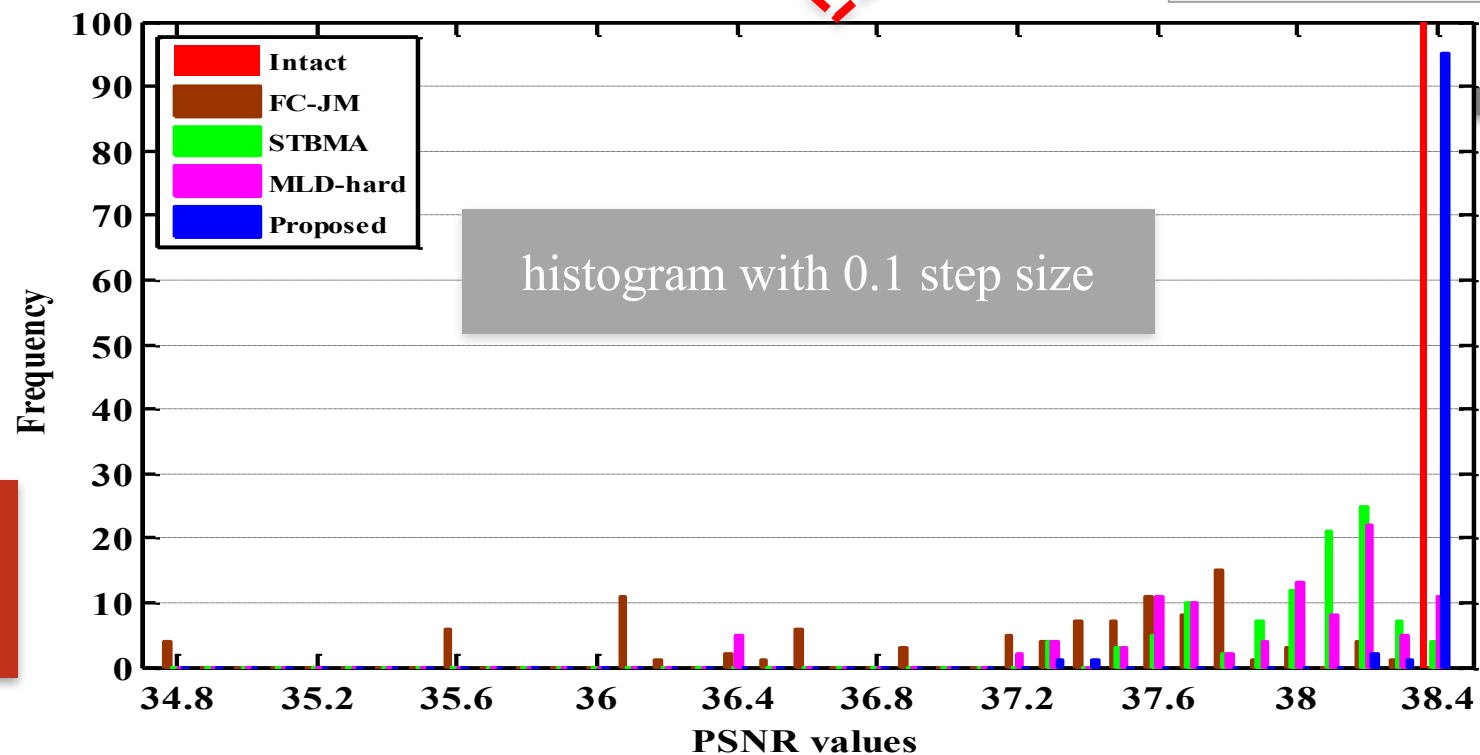
## ❖ Decoding Options :

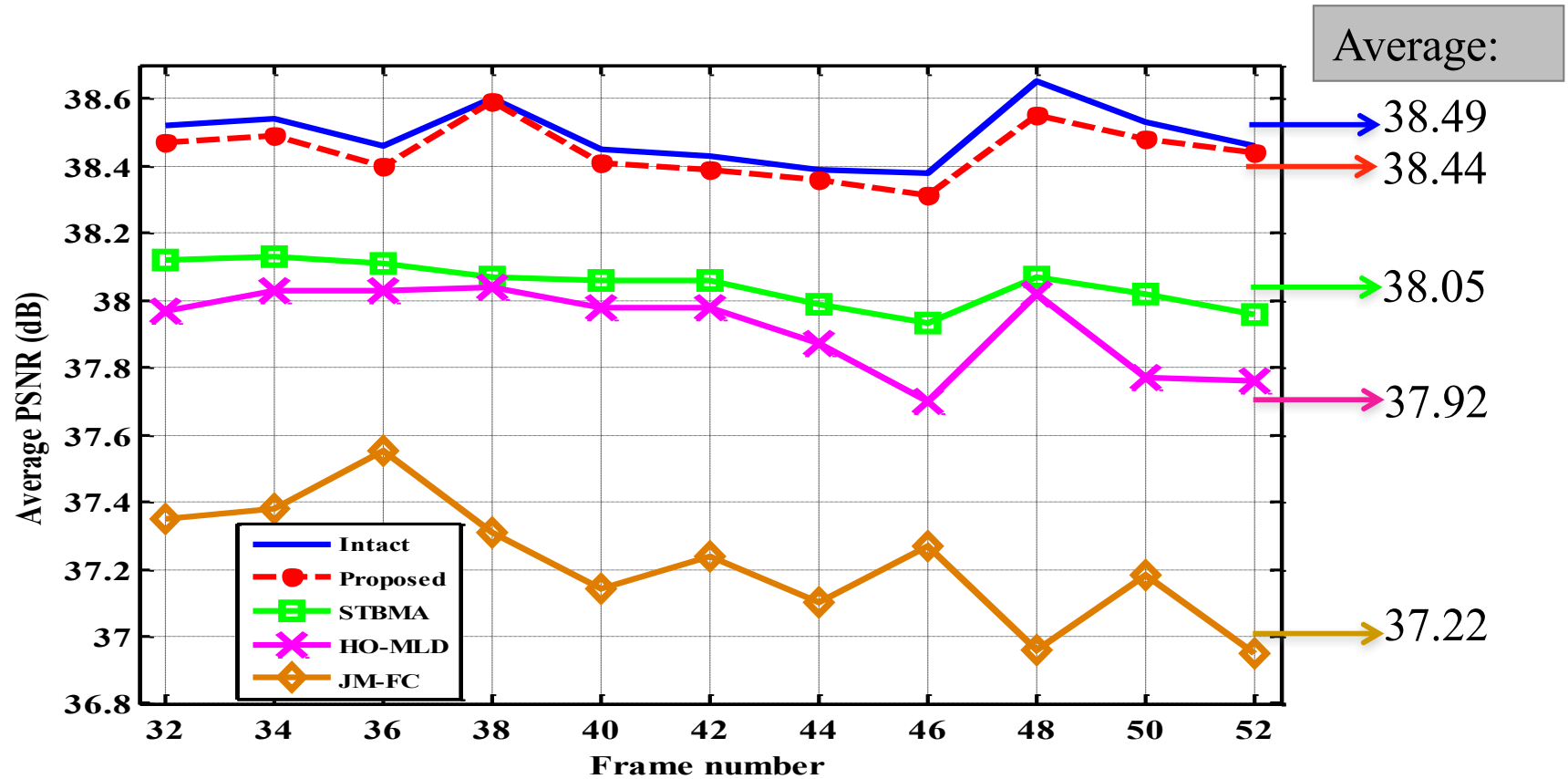
- Error Concealment:** Frame Copy (**FC-JM**)  
Spatio Temporal Boundary Matching (**STBMA**)
- Error Correction:** Maximum likelihood Decoding (**HO-MLD**)  
Proposed List Decoding (**Proposed**)

Seq: Crew, QP 27



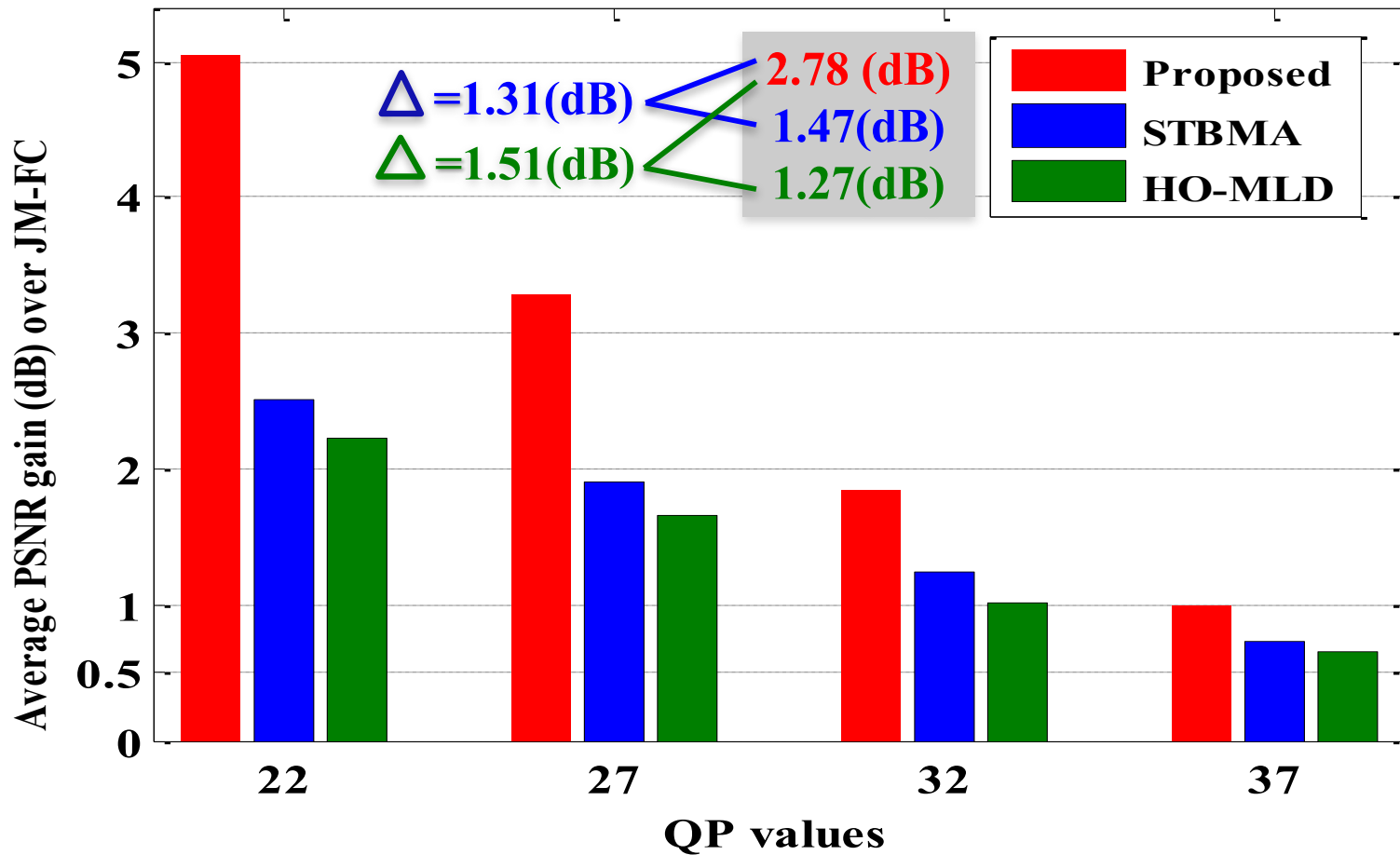
95% very close PSNR to the intact one.







## PSNR gain over JM-FC



➤ Percentage of fully correcting bitstream:

**Proposed approach: 80%**

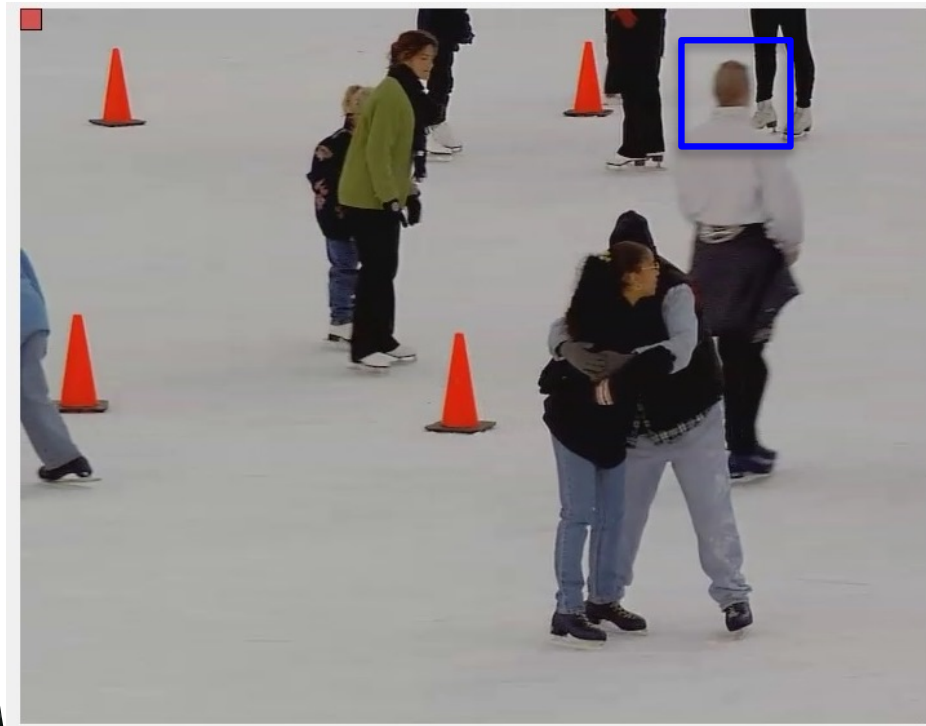
**HO-MLD approach: 6%**

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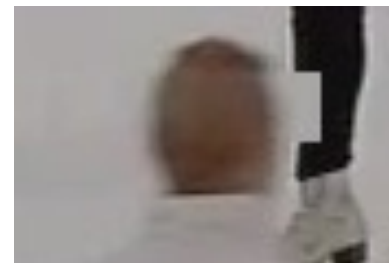
# Visual quality comparison



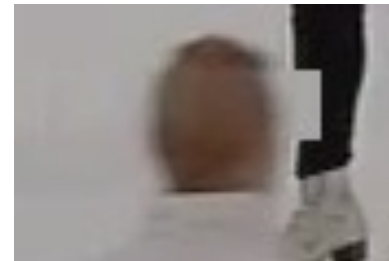
Intact



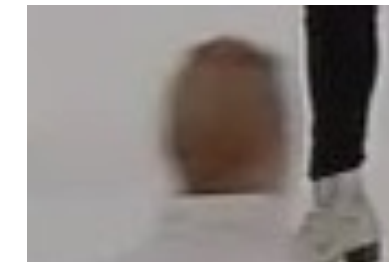
JM-FC



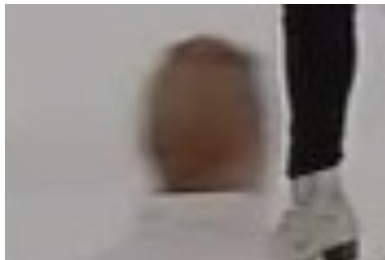
STBMA



HO-MLD



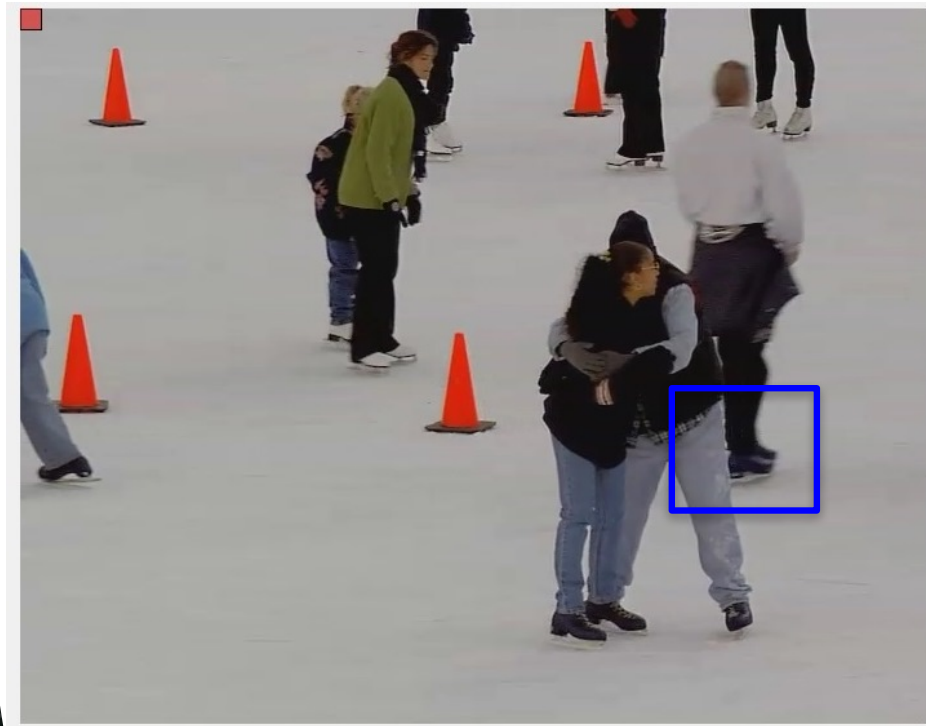
Proposed



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# Visual quality comparison



Intact



JM-FC



STBMA



HO-MLD



Proposed



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# Conclusions



## One bit in error:

- Candidate reduction from  $N$  to  $N/32$ ;
- Complexity reduction : **97%**

## On average:

- **2.78 dB** gains over **JM-FC**
- **80%** fully corrected packet

## □ Future step:

- The proposed approach can also be applied on HEVC coded sequences, and it can be extended to the case of more than one bit in error [1].



Thank you  
for your attention!



Questions

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